

To test the heuristic functions  $h_1$  and  $h_2$ , we generated 1200 random problems with solution lengths from 2 to 24 (100 for each even number) and solved them with iterative deepening search and with A\* tree search using both  $h_1$  and  $h_2$ . Figure 3.29 gives the average number of nodes generated by each strategy and the effective branching factor. The results suggest that  $h_2$  is better than  $h_1$ , and is far better than using iterative deepening search. Even for small problems with  $d = 12$ , A\* with  $h_2$  is 50,000 times more efficient than uninformed iterative deepening search.

$d$	Search Cost (nodes generated)			Effective Branching Factor		
	IDS	A*( $h_1$ )	A*( $h_2$ )	IDS	A*( $h_1$ )	A*( $h_2$ )
2	10	6	6	2.45	1.79	1.79
4	112	13	12	2.87	1.48	1.45
6	680	20	18	2.73	1.34	1.30
8	6384	39	25	2.80	1.33	1.24
10	47127	93	39	2.79	1.38	1.22
12	3644035	227	73	2.78	1.42	1.24
14	–	539	113	–	1.44	1.23
16	–	1301	211	–	1.45	1.25
18	–	3056	363	–	1.46	1.26
20	–	7276	676	–	1.47	1.27
22	–	18094	1219	–	1.48	1.28
24	–	39135	1641	–	1.48	1.26

**Figure 3.29** Comparison of the search costs and effective branching factors for the ITERATIVE-DEEPENING-SEARCH and A\* algorithms with  $h_1$ ,  $h_2$ . Data are averaged over 100 instances of the 8-puzzle for each of various solution lengths  $d$ .

#### DOMINATION

One might ask whether  $h_2$  is *always* better than  $h_1$ . The answer is “Essentially, yes.” It is easy to see from the definitions of the two heuristics that, for any node  $n$ ,  $h_2(n) \geq h_1(n)$ . We thus say that  $h_2$  **dominates**  $h_1$ . Domination translates directly into efficiency: A\* using  $h_2$  will never expand more nodes than A\* using  $h_1$  (except possibly for some nodes with  $f(n) = C^*$ ). The argument is simple. Recall the observation on page 97 that every node with  $f(n) < C^*$  will surely be expanded. This is the same as saying that every node with  $h(n) < C^* - g(n)$  will surely be expanded. But because  $h_2$  is at least as big as  $h_1$  for all nodes, every node that is surely expanded by A\* search with  $h_2$  will also surely be expanded with  $h_1$ , and  $h_1$  might cause other nodes to be expanded as well. Hence, it is generally better to use a heuristic function with higher values, provided it is consistent and that the computation time for the heuristic is not too long.

### 3.6.2 Generating admissible heuristics from relaxed problems

We have seen that both  $h_1$  (misplaced tiles) and  $h_2$  (Manhattan distance) are fairly good heuristics for the 8-puzzle and that  $h_2$  is better. How might one have come up with  $h_2$ ? Is it possible for a computer to invent such a heuristic mechanically?

$h_1$  and  $h_2$  are estimates of the remaining path length for the 8-puzzle, but they are also perfectly accurate path lengths for *simplified* versions of the puzzle. If the rules of the puzzle